

## ON BLOCK MATRICES ASSOCIATED WITH DISCRETE TRIGONOMETRIC TRANSFORMS AND THEIR USE IN THE THEORY OF WAVE PROPAGATION\*

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### Abstract

Block matrices associated with discrete Trigonometric transforms (DTT's) arise in the mathematical modelling of several applications of wave propagation theory including discretizations of scatterers and radiators with the Method of Moments, the Boundary Element Method, and the Method of Auxiliary Sources. The DTT's are represented by the Fourier, Hartley, Cosine, and Sine matrices, which are unitary and offer simultaneous diagonalizations of specific matrix algebras. The main tool for the investigation of the aforementioned wave applications is the efficient inversion of such types of block matrices. To this direction, in this paper we develop an efficient algorithm for the inversion of matrices with  $U$ -diagonalizable blocks ( $U$  a fixed unitary matrix) by utilizing the  $U$ -diagonalization of each block and subsequently a similarity transformation procedure. We determine the developed method's computational complexity and point out its high efficiency compared to standard inversion techniques. An implementation of the algorithm in Matlab is given. Several numerical results are presented demonstrating the CPU-time efficiency and accuracy for ill-conditioned matrices of the method. The investigated matrices stem from real-world wave propagation applications.

*Mathematics subject classification:* 65F05, 65T50, 74J20, 78A40, 15A09.

*Key words:* Discrete Trigonometric transforms, Block matrices, Efficient inversion algorithms, Wave radiation and scattering, Numerical methods in wave propagation theory.

### 1. Introduction

Discrete Trigonometric transforms (DTT's) play a significant role in wave scattering and radiation theory, signal processing, physics, and numerical linear algebra. Representative examples constitute the discrete Fourier transforms (DFT's), the discrete Hartley transforms (DHT's), the discrete Cosine transforms (DCT's), and the discrete Sine transforms (DST's). Their primary contribution lies in the significant reduction of the complexity in the associated mathematical problems. For example, applications of such appropriate transforms in differential and integral equations reduce them to algebraic equations, whose solutions are more easily obtained, see, e.g., [1, 2]. Moreover, in harmonic analysis as well as in signal processing the DFT decomposes a signal sequence into its frequency components [3]. It is important to note that this wide applicability of the DTT's is mainly justified by the existence of fast algorithms, that allow the transforms computations within  $\mathcal{O}(n \log_2 n)$  (instead of  $\mathcal{O}(n^2)$ ) when performing directly the matrix-vector product of length  $n$ ) [4–6].

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In particular, block matrices associated with DTT's arise in concrete physical and technological applications including: (i) the solution of wave scattering and radiation problems with the Method of Moments (MoM) [7, 8], (ii) the investigation and optimization of numerical methods for electromagnetic scattering problems, such as the Method of Auxiliary Sources (MAS) [9]- [10], (iii) the numerical solution of integral equations with the Boundary Element Method (BEM) [11], (iv) optical imaging [12], (v) image compression [13], (vi) efficient preconditioning of Toeplitz systems [14]. Besides, we point out that these applications exhibit the essential role of the inversion of such types of complex block matrices for the derivation of formulas determining the error bounds of numerical methods as well as for the numerical or semi-analytical computation of solutions [15].

The specific DTT's mentioned above are represented by the Fourier, Hartley, Cosine, and Sine matrices. These unitary matrices offer simultaneous diagonalizations of specific matrix algebras, including circulants, skew-circulants, Toeplitz-plus-Hankel, tridiagonal [16]. These matrix algebras are unified by considering the algebra  $\text{Diag}(U)$  of all  $U$ -diagonalizable matrices, for  $U$  a fixed unitary matrix.

For the mathematical modelling of the above mentioned wave applications we develop in this paper an efficient method for the inversion of an  $m \times m$  block matrix  $A = [A_{ij}]$  with  $U$ -diagonalizable blocks of order  $n$  ( $i, j = 1, \dots, m$ ). First, we consider the diagonalizations  $A_{ij} = U\Lambda_{ij}U^*$ , where  $\Lambda_{ij}$  is the diagonal matrix containing the eigenvalues of  $A_{ij}$ , and hence the inversion of  $A$  is reduced to that of the  $m \times m$  block matrix  $\Lambda = [\Lambda_{ij}]$  with diagonal blocks of order  $n$ . For the inversion of  $\Lambda$  we construct, by using concepts of Graph Theory, an appropriate permutation matrix  $P$  so that the matrix  $P\Lambda P^T = \text{diag}(\Lambda'_1, \Lambda'_2, \dots, \Lambda'_n)$  is block-diagonal with  $\Lambda'_k$  invertible  $m \times m$  full matrices. The inverse  $\Lambda^{-1} = [L_{ij}]$  of  $\Lambda$  is then determined by inverting each block  $\Lambda'_k$  with a standard LU direct solver. Finally, the inverse of  $A$  is given by the inverse block-diagonalization that is  $A^{-1} = [UL_{ij}U^*]$ . An implementation in Matlab of the above described algorithmic inversion is given in the Appendix.

For any one of the choices of  $U$ , that is Fourier, Hartley, Cosine, and Sine matrices, the matrix multiplications  $U^*A_{ij}U$  and  $UL_{ij}U^*$ , appearing in the block-diagonalizations, are computed by applying the DTT's for  $n \neq 2^p$  or the fast Trigonometric transforms (FTT's) for  $n=2^p$ , that is the fast Fourier, Hartley, Cosine, and Sine transforms [4]- [6]. Hence, the computational complexity (i.e. the total number of required scalar complex multiplications) of the inversion algorithm is  $n\mathcal{O}(m^3) + 2m^2n^2$  for  $n \neq 2^p$  and  $n\mathcal{O}(m^3) + 2m^2\mathcal{O}(n\log_2 n)$  for  $n=2^p$ . This shows that the developed method is far more efficient than the LU decomposition applied to the original matrix  $A$ , having complexity  $\mathcal{O}(m^3n^3)$ . We note that in several wave applications the order  $n$  of each block may be chosen equal to  $2^p$  by selecting suitable discretizations of the scatterer's or radiator's surface [7]- [10].

On the other hand, the above described inversion method can be also applied for the efficient determination of the eigenvalues of a matrix  $A$  with  $U$ -diagonalizable blocks. Specifically, the eigenvalues of  $A$  are those of all blocks  $\Lambda'_k$  and thus their computation requires  $n\mathcal{O}(m^3) + m^2\mathcal{O}(n\log_2 n)$  multiplications. Besides, we notice the parallel nature of the proposed inversion algorithm, since the inversion of each specific block  $\Lambda'_k$  can be handled by a different processing unit. For a discussion on parallel algorithms for inverting block matrices which arise in inverse wave scattering theory see [17].

Several numerical results are presented exhibiting the efficiency of the proposed method and highlighting its beneficial contribution in the numerical implementation of certain scattering and radiation applications. We compare in terms of CPU time the developed algorithmic

inversion with the LU method, applied to the initial matrix  $A$ , and investigate the achieved high speedup for various matrices with circulant, skew-circulant, Toeplitz-plus-Hankel, and tridiagonal blocks, appearing in real-world radiation applications, and in particular in the modelling of radiation by circular-loop antennas with the MoM [7]- [8]. Moreover, we examine the accuracy of the proposed inversion method for the case of ill-conditioned matrices, arising in the numerical solution of electromagnetic scattering problems by layered objects with the MAS [9]- [10]. The achieved inversion error is significantly reduced by the application of the developed method compared with that of the LU decomposition. This fact is justified numerically by the decrease of the conditions numbers of the blocks  $\Lambda'_k$ , which have to be inverted by the developed algorithm, and a related discussion is included.

## 2. Preliminaries, Notations and Terminology

A *Discrete Trigonometric Transform* (DTT) of length  $n$  is a linear transformation

$$T_M : \mathbb{C}^n \rightarrow \mathbb{C}^n, \quad T_M \mathbf{x} = M \mathbf{x}, \tag{2.1}$$

where  $M$  is a Fourier, Hartley, Cosine, or Sine square matrix of order  $n$ . The respective DTT is named the *Discrete Fourier, Hartley, Cosine, or Sine Transform* (DFT, DHT, DCT, or DST) of length  $n$ . We collect the 4 types I, II, III, and IV of each one of the  $n \times n$  Fourier, Hartley, Cosine, and Sine matrices, commonly used in the literature (see, e.g., [6, 16, 18]) by  $(j, k \in \{0, \dots, n - 1\})$

$$[F_n^I]_{jk} = \frac{1}{\sqrt{n}} w_n^{jk}, \quad [F_n^{II}]_{jk} = \frac{1}{\sqrt{n}} w_{2n}^{j(2k+1)}, \tag{2.2a}$$

$$[F_n^{III}]_{jk} = \frac{1}{\sqrt{n}} w_{2n}^{(2j+1)k}, \quad [F_n^{IV}]_{jk} = \frac{1}{\sqrt{n}} w_{4n}^{(2j+1)(2k+1)}, \tag{2.2b}$$

$$[H_n^I]_{jk} = \frac{1}{\sqrt{n}} \text{cas} \left( \frac{2\pi jk}{n} \right), \quad [H_n^{II}]_{jk} = \frac{1}{\sqrt{n}} \text{cas} \left( \frac{2\pi j(2k+1)}{2n} \right), \tag{2.3a}$$

$$[H_n^{III}]_{jk} = \frac{1}{\sqrt{n}} \text{cas} \left( \frac{2\pi(2j+1)k}{2n} \right), \quad [H_n^{IV}]_{jk} = \frac{1}{\sqrt{n}} \text{cas} \left( \frac{2\pi(2j+1)(2k+1)}{4n} \right), \tag{2.3b}$$

$$[C_n^I]_{jk} = \sqrt{\frac{2}{n-1}} \delta_j \delta_k \epsilon_j \epsilon_k \cos \left( \frac{jk\pi}{n-1} \right), \quad [C_n^{II}]_{jk} = \sqrt{\frac{2}{n}} \delta_j \cos \left( \frac{j(k+\frac{1}{2})\pi}{n} \right),$$

$$[C_n^{III}]_{jk} = \sqrt{\frac{2}{n}} \delta_k \cos \left( \frac{(j+\frac{1}{2})k\pi}{n} \right), \quad [C_n^{IV}]_{jk} = \sqrt{\frac{2}{n}} \cos \left( \frac{(j+\frac{1}{2})(k+\frac{1}{2})\pi}{n} \right), \tag{2.4}$$

$$[S_n^I]_{jk} = \sqrt{\frac{2}{n+1}} \sin \left( \frac{(j+1)(k+1)\pi}{n+1} \right), \quad [S_n^{II}]_{jk} = \sqrt{\frac{2}{n}} \epsilon_j \sin \left( \frac{(j+1)(k+\frac{1}{2})\pi}{n} \right),$$

$$[S_n^{III}]_{jk} = \sqrt{\frac{2}{n}} \epsilon_k \sin \left( \frac{(j+\frac{1}{2})(k+1)\pi}{n} \right), \quad [S_n^{IV}]_{jk} = \sqrt{\frac{2}{n}} \sin \left( \frac{(j+\frac{1}{2})(k+\frac{1}{2})\pi}{n} \right). \tag{2.5}$$

where  $w_n = \exp(-2\pi i/n)$ ,  $i^2 = -1$ ,

$$\text{cas}(t) := \cos(t) + \sin(t),$$



Table 2.1: The coefficient vectors  $\beta_X^c$  and  $\beta_X^s$  for each type of the DCT's and DST's.

	$\beta_X^c$
DCT-I	$(0, \sqrt{2}, \sqrt{2}, 0, \sqrt{2}, \sqrt{2})$
DCT-II	$(1, 1, 1, 1, 1, 1)$
DCT-III	$(0, \sqrt{2}, 1, 0, \sqrt{2}, 1)$
DCT-IV	$(1, 1, 1, -1, 1, 1)$
	$\beta_X^s$
DST-I	$(0, 1, 1, 0, 1, 1)$
DST-II	$(-1, 1, 1, -1, 1, 1)$
DST-III	$(0, 1, \sqrt{2}, 0, 1, \sqrt{2})$
DST-IV	$(-1, 1, 1, 1, 1, 1)$

Now, under the above considerations and summarizing results of [16, 18], we have

$$\begin{aligned}
 \text{Diag}(F_n^I) &= \text{Diag}(F_n^{III}) = \mathcal{C}_n, & \text{Diag}(F_n^{II}) &= \text{Diag}(F_n^{IV}) = \mathcal{S}\mathcal{C}_n, \\
 \text{Diag}(H_n^I) &= \mathcal{C}_n^0 \oplus J'_n \mathcal{C}_n^1, & \text{Diag}(H_n^{II}) &= \mathcal{S}\mathcal{C}_n^1 \oplus J''_n \mathcal{S}\mathcal{C}_n^0, \\
 \text{Diag}(H_n^{III}) &= \mathcal{C}_n^0 \oplus J_n \mathcal{C}_n^1, & \text{Diag}(H_n^{IV}) &= \mathcal{S}\mathcal{C}_n^1 \oplus J_n \mathcal{S}\mathcal{C}_n^0, \\
 \text{Diag}(C_n^X) &= \mathcal{B}_n(\beta_X^c), & \text{Diag}(S_n^X) &= \mathcal{B}_n(\beta_X^s),
 \end{aligned} \tag{2.8}$$

where  $\oplus$  denotes the *orthogonal sum of linear subspaces* of  $\mathbb{C}^{n \times n}$  with respect to the Frobenius inner product [16]. In particular, from (2.8),  $\text{Diag}(H_n^X)$  are subspaces of  $\mathcal{T}_n + \mathcal{H}\mathcal{A}_n$ , and hence every matrix of  $\text{Diag}(H_n^X)$  may be expressed as a symmetric Toeplitz plus Hankel matrix.

Besides, concrete wave propagation applications, analyzed in [7]- [10], demonstrate the significance of the generalization of the notion of *U*-diagonalizable matrices to that of *matrices with U-diagonalizable blocks*, i.e.  $m \times m$  block matrices  $A = [A_{ij}]$  with  $A_{ij}$   $n \times n$  *U*-diagonalizable matrices.

### 3. Inversion Algorithm

We develop an algorithmic approach for the inversion of an  $m \times m$  block matrix  $A = [A_{ij}]$ , with *U*-diagonalizable blocks of order  $n$  based on: (i) the block-diagonalization of  $A$  and (ii) a similarity transformation algorithm for the inversion of the block matrix, determined by the eigenvalues of the blocks  $A_{ij}$ .

For the development of the inversion algorithm we utilize essentially the following

**Theorem 3.1.** *Let  $U$  be a unitary matrix of order  $n$  and  $A = [A_{ij}]$  an  $m \times m$  block matrix with *U*-diagonalizable blocks of order  $n$ . Then,  $A$  is block-diagonalized as*

$$A = \text{diag}(U, U, \dots, U) [\Lambda_{ij}] \text{diag}(U^*, U^*, \dots, U^*), \quad \Lambda_{ij} = \text{diag}(\lambda_{ij}^1, \lambda_{ij}^2, \dots, \lambda_{ij}^n), \tag{3.1}$$

where  $\lambda_{ij}^k$  are the eigenvalues of  $A_{ij}$  ( $k \in \{1, \dots, n\}$  and  $i, j \in \{1, \dots, m\}$ ) and all  $\Lambda_{ij}$  have as eigenvectors the columns of  $U$ .

Besides, if  $A$  is invertible, the inverse  $A^{-1}$  has the block-diagonalization

$$A^{-1} = \text{diag}(U, U, \dots, U) [L_{ij}] \text{diag}(U^*, U^*, \dots, U^*), \quad \Lambda^{-1} = [L_{ij}]. \tag{3.2}$$

Eq. (3.2) reduces the inversion of  $A$  to that of the  $m \times m$  block matrix  $\Lambda = [\Lambda_{ij}]$  with diagonal blocks of order  $n$ . The inverse  $\Lambda^{-1}$  is computed below by applying a similarity transformation algorithm.

A similar block-diagonalization to that of the previous Theorem has also been considered in [21] for the particular case of matrices with circulant blocks (each diagonalized by the Fourier matrix). However, the inversion of the respective matrix  $\Lambda$  in [21] is obtained by a recursive method, different than the one presented hereafter.

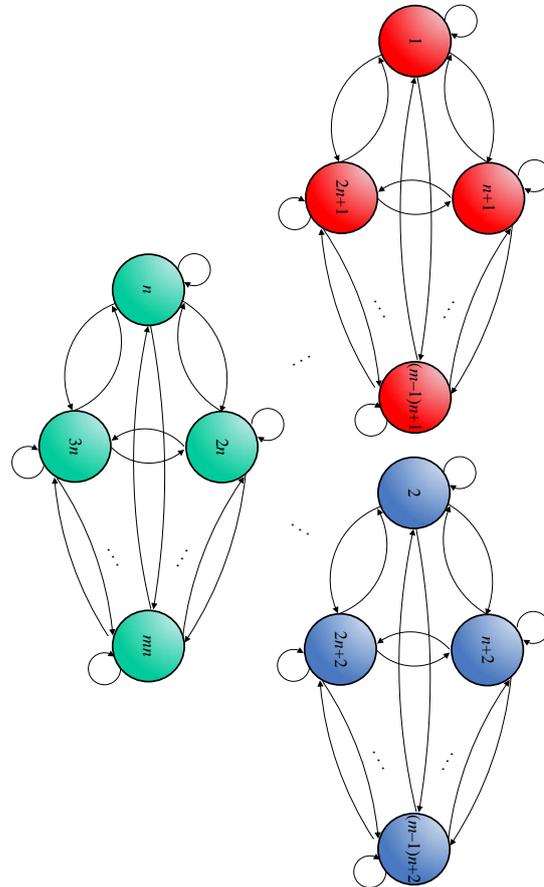


Fig. 3.1. Graph  $G(\Lambda)$  of the  $m \times m$  block matrix  $\Lambda = [\Lambda_{ij}]$  with diagonal blocks of order  $n$

The basic idea for the inversion of  $\Lambda$  lies in the application of a suitable similarity transformation, making the matrix with diagonal blocks  $\Lambda$  similar to a block-diagonal matrix  $\Lambda'$ . To this direction, we construct the graph  $G(\Lambda)$  of  $\Lambda$  (see Fig. 3.1), where only the diagonal elements of the blocks  $\Lambda_{ij}$  are considered, including the zero ones. In view of Fig. 3.1, the graph  $G(\Lambda)$  is a union of  $n$  disjoint subgraphs each with  $m$  nodes. Thus, taking into account this decomposition of the graph, we construct the permutation matrix

$$P = [\mathbf{e}_1, \mathbf{e}_{n+1}, \dots, \mathbf{e}_{(m-1)n+1}, \mathbf{e}_2, \mathbf{e}_{n+2}, \dots, \mathbf{e}_{(m-1)n+2}, \dots, \mathbf{e}_n, \mathbf{e}_{2n}, \dots, \mathbf{e}_{mn}]^T,$$

where  $(\mathbf{e}_i)_{i \in \{1, \dots, mn\}}$  the standard base of  $\mathbb{R}^{mn}$ . By the definition of  $P$  we get the similarity transformation

$$P\Lambda P^T = \Lambda'_1 \oplus \Lambda'_2 \oplus \dots \oplus \Lambda'_n \equiv \Lambda', \tag{3.3}$$

with

$$[\Lambda'_k]_{ij} = \lambda_{ij}^k \quad (k \in \{1, \dots, n\}, \quad i, j \in \{1, \dots, m\}).$$

Thus, so far the original  $m \times m$  block matrix  $\Lambda$  with diagonal blocks  $\Lambda_{ij}$  of order  $n$  has been transformed to the block-diagonal matrix  $\Lambda'$  with  $n$  full blocks  $\Lambda'_k$  of order  $m$ .

Furthermore, the inverse  $\Lambda'^{-1}$  of  $\Lambda'$  is determined by inverting each block  $\Lambda'_k$  with a standard direct solver such as the LU decomposition. We note that the blocks  $\Lambda'_k$  are invertible in case the matrix  $\Lambda$  is invertible and the inverse  $\Lambda'^{-1}$  is given by

$$\Lambda'^{-1} = \Lambda_1'^{-1} \oplus \Lambda_2'^{-1} \oplus \dots \oplus \Lambda_n'^{-1}, \quad [\Lambda_k'^{-1}]_{ij} = \ell_{ij}^k. \tag{3.4}$$

Now,  $\Lambda^{-1} = [L_{ij}]$  is derived from  $\Lambda'^{-1}$  by applying the inverse similarity transformation

$$\Lambda^{-1} = P^T \Lambda'^{-1} P, \tag{3.5}$$

which implies that the  $n \times n$  diagonal blocks  $L_{ij}$  of the matrix  $\Lambda^{-1}$  are of the form

$$L_{ij} = \text{diag}(\ell_{ij}^1, \ell_{ij}^2, \dots, \ell_{ij}^n). \tag{3.6}$$

Finally, application of the block-diagonalization (3.2) to  $\Lambda^{-1}$  gives the inverse  $A^{-1}$  of  $A$ .

### 4. Computational Complexity of the Inversion

The computational complexity, which we consider, expresses the total number of required scalar complex multiplications. For the determination of the computational complexity  $C(m,n)$  of the proposed algorithm for the inversion of a matrix  $A$  with  $U$ -diagonalizable blocks we need the complexities: (i)  $C_{bdA}$  of the block-diagonalization of  $A$ , (ii)  $C_S$  of the similarity transformation algorithm for the inversion of  $\Lambda$ , and (iii)  $C_{bdA^{-1}}$  of the block-diagonalization of  $A^{-1}$ . Clearly, the complexity  $C(m,n)$  is the sum

$$C(m,n) = C_{bdA}(m,n) + C_S(m,n) + C_{bdA^{-1}}(m,n). \tag{4.1}$$

The unitary matrix  $U$  is introduced to unify the 16 cases of the Fourier, Hartley, Cosine, and Sine matrices, described in (2.2)-(2.5). For these choices of  $U$  the matrix multiplications  $U^* A_{ij} U$  and  $U L_{ij} U^*$ , appearing in the block-diagonalizations (3.1) and (3.2), may be computed by applying the discrete (DTT's) or the fast (FTT's) Trigonometric transforms.

The following Theorem gives the complexity  $C(m,n)$  of the algorithmic inversion

**Theorem 4.1.** *Let  $A = [A_{ij}]$  be an invertible  $m \times m$  block matrix with  $U$ -diagonalizable blocks of order  $n$ , where  $U$  represents any one of  $F_n^X, H_n^X, C_n^X$ , or  $S_n^X$  for  $X \in \{I, II, III, IV\}$ . Then, the computational complexity  $C(m,n)$  of the algorithmic inversion of  $A$ , developed in Section 3, is given by*

$$C(m,n) = \begin{cases} n\mathcal{O}(m^3) + 2m^2n^2, & \log_2 n \notin \mathbb{N}, \\ n\mathcal{O}(m^3) + 2m^2\mathcal{O}(n\log_2 n), & \log_2 n \in \mathbb{N}. \end{cases} \tag{4.2}$$

*Proof.* The complexity  $C_{bdA}$  is that of the determination of the blocks  $\Lambda_{ij}$  of (3.1), which are computed for  $n \neq 2^p$  by the DTT's (2.1)-(2.5), requiring  $n^2$  multiplications, and for  $n=2^p$  by the FTT's, requiring  $\mathcal{O}(n\log_2 n)$  multiplications (see [5], [6]), yielding

$$C_{bdA}(m,n) = \begin{cases} m^2n^2, & \log_2 n \notin \mathbb{N}, \\ m^2\mathcal{O}(n\log_2 n), & \log_2 n \in \mathbb{N}. \end{cases}$$

Furthermore,  $C_S$  is determined by the complexity of the inversions of the  $n$  blocks  $\Lambda'_k$  of order  $m$  of matrix  $\Lambda'$  in (3.3) with an LU solver, resulting  $C_S(m, n) = n\mathcal{O}(m^3)$ . Note that the similarity transformations (3.3) and (3.5) do not contribute to  $C_S$ , since they involve only rows and columns interchanges.

Finally, in order to compute the block-diagonalization (3.2), we apply the inverse DTT's or FTT's to the blocks  $L_{ij}$  of (3.6). Since both the DTT's and the inverse DTT's require  $n^2$  multiplications and the respective FTT's  $\mathcal{O}(n\log_2 n)$  multiplications, we conclude that  $C_{bdA^{-1}}(m, n) = C_{bdA}(m, n)$ . Thus, the desired (4.2) follows by (4.1).  $\square$

The following remarks concern useful aspects of the inversion's computational complexity

- (i) *Efficient computation of the eigenvalues of matrices with U-diagonalizable blocks:* by combining (3.1) and (3.3) we have

$$A = (U \oplus U \oplus \dots \oplus U) P^T (\Lambda'_1 \oplus \Lambda'_2 \oplus \dots \oplus \Lambda'_n) P (U^* \oplus U^* \oplus \dots \oplus U^*).$$

Thus, the eigenvalues of  $A$  coincide with those of all blocks  $\Lambda'_k$  and their complexity is

$$C_{eig}(m, n) = \begin{cases} n\mathcal{O}(m^3) + m^2n^2, & \log_2 n \notin \mathbb{N}, \\ n\mathcal{O}(m^3) + m^2\mathcal{O}(n\log_2 n), & \log_2 n \in \mathbb{N}. \end{cases} \tag{4.3}$$

The factor 2 in  $C$  of (4.2) does not appear in  $C_{eig}$  of (4.3), since for the eigenvalues computation are required only the direct (and not the inverse) DTT's or FTT's.

- (ii) *Parallel nature of the proposed algorithm:* for  $n$  large and  $m$  small enough, the block-diagonalizations (3.1) and (3.2) require  $m^2$  processing nodes for the execution of each one of the direct and inverse DTT's, and then the complexity is decreased to

$$C_{par1}(m, n) = n\mathcal{O}(m^3) + 2\mathcal{O}(n\log_2 n).$$

Moreover, for  $m$  large and  $n$  small enough, the inversion of each block  $\Lambda'_k$  may be assigned to a different one from  $n$  processing nodes and the complexity reduces to

$$C_{par2}(m, n) = \mathcal{O}(m^3) + 2m^2\mathcal{O}(n\log_2 n).$$

- (iii) From (4.2), the algorithmic inversion works best for  $n = 2^p$ , where the block-diagonalizations are computed at "FTT-speed". On the other hand, the complexity of the LU decomposition and the numerical inversion of the matrix  $A$  is of order  $C_{LU}(m, n) = \mathcal{O}(m^3n^3)$  and hence the proposed algorithmic inversion is much faster than the LU method (as shown also in the numerical results of Section 5).

- (iv) The accurate numerical solutions of certain electromagnetic scattering and radiation problems [7–10] require large  $n$  and small  $m$  (usually 2 or 4), yielding

$$C(n) = \begin{cases} \mathcal{O}(n^2), & \log_2 n \notin \mathbb{N} \\ \mathcal{O}(n\log_2 n), & \log_2 n \in \mathbb{N} \end{cases}, \quad C_{LU}(n) = \mathcal{O}(n^3),$$

and thus  $C_{LU}(n)$  is at least one order of magnitude greater than  $C(n)$ .

## 5. Numerical Results

In this Section we present numerical results for the developed inversion algorithm with respect to (i) the execution (CPU) time and (ii) the accuracy for ill-conditioned matrices. The investigated block matrices stem from real-world scattering and radiation applications. The algorithm is implemented by the Matlab function *inv\_unit\_diagon\_blocks*, given in the Appendix. The numerical results are compared with the LU decomposition method, using Gaussian elimination with partial pivoting, and the efficiency of the developed algorithm is discussed. All numerical experiments are implemented in Matlab 7 with double precision arithmetic and executed in an Intel Pentium M processor 1.60 GHz with 504 MB of RAM.

### 5.1. Execution Time

First, we consider matrices with circulant and skew-circulant blocks, appearing in the real-world application of the modelling of radiation by circular-loop antenna arrays with the Method of Moments [7, 8]. For this application  $m$  represents the number of circular loops, while  $n$  the number of basis functions utilized. Figs. 5.1a and 5.1b depict the CPU time for the inversions with respect to  $mn$  of the present method and the LU decomposition for  $m \times m$  block matrices with  $F_n^I$ - or  $F_n^{III}$ -diagonalizable circulant blocks of order  $n$  for fixed (a)  $n=100, 200, 400$  and (b)  $m=4, 8, 16$ . The situation is similar in the case of block matrices with  $F_n^{II}$  or  $F_n^{IV}$ -diagonalizable skew-circulant blocks. The developed algorithmic inversion is clearly faster than the LU inversion and the logarithmic scale in Fig. 5.1 emphasizes its lower complexity. This CPU-time improvement, achieved by the developed method, may lead to the efficient and fast modelling of electrically large circular-loop antenna arrays.

Furthermore, we consider matrices with symmetric Toeplitz-plus-Hankel blocks, composed of random real entries, chosen from a normal distribution with mean zero, variance one, and standard deviation one. The Toeplitz and the Hankel counterparts of these blocks, belonging to  $H_n^X$  matrix algebras, are constructed by following the general guidelines of [16], modified into the present context in Table 5.1.

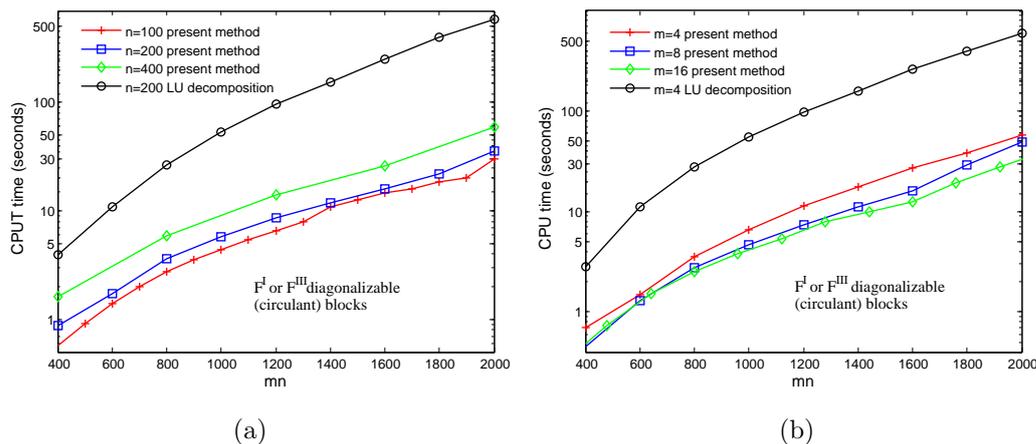


Fig. 5.1. CPU time for the inversions with respect to  $mn$  of the present method and the LU decomposition for  $m \times m$  block matrices with  $F_n^I$ - or  $F_n^{III}$ -diagonalizable circulant blocks of order  $n$  for: (a)  $n=100, 200, 400$  and (b)  $m=4, 8, 16$ , appearing in modelling of radiation by circular-loop antenna arrays.

Figs. 5.2a and 5.2b show the CPU time for the inversions with respect to  $mn$  of the present method and the LU decomposition for  $m \times m$  block matrices with (a)  $H_n^I$ - and (b)  $H_n^{III}$ -diagonalizable Toeplitz-plus-Hankel blocks of order  $n$  for  $n=300, 400, 500$ . By the statements of Fig. 5.2 the present algorithmic inversion is carried out in a few seconds even for large  $mn$ . However, the computational cost of the LU inversion increases rapidly with  $mn$  and this fact makes the application of this inversion prohibitive. Moreover, for arbitrary fixed  $n$  ( $m$ ) the difference between the CPU time of the LU and the present inversion method increases significantly with  $m$  ( $n$ ). Similar conclusions hold for block matrices with  $H_n^{II}$  or  $H_n^{IV}$ -diagonalizable Toeplitz-plus-Hankel blocks.

Finally, we consider tridiagonal matrices of the form (2.7). Figs. 5.3a and 5.3b depict the inversion CPU time with respect to  $n$  of the present method and the LU decomposition for (a)  $C_n^X$ - and (b)  $S_n^X$ -diagonalizable matrices for  $X=I, II, III, IV$ . For large  $n$  the achieved speedup of the developed method compared to the LU decomposition increases rapidly. For fixed  $n$  the CPU time of the algorithmic inversion does not vary significantly with  $X$ .

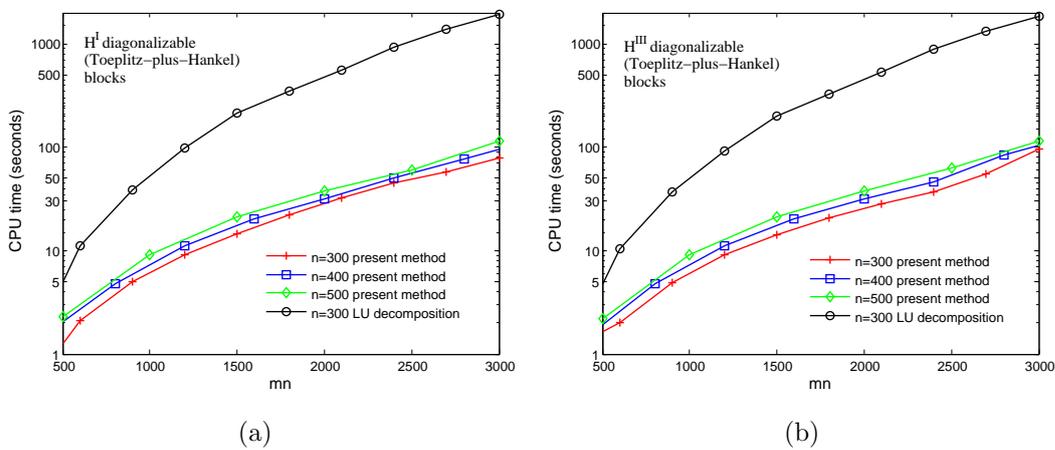


Fig. 5.2. CPU time for the inversions with respect to  $mn$  of the present method and the LU decomposition for  $m \times m$  block matrices with (a)  $H_n^I$ - and (b)  $H_n^{III}$ -diagonalizable Toeplitz-plus-Hankel blocks of order  $n$  for  $n=300, 400, 500$ .

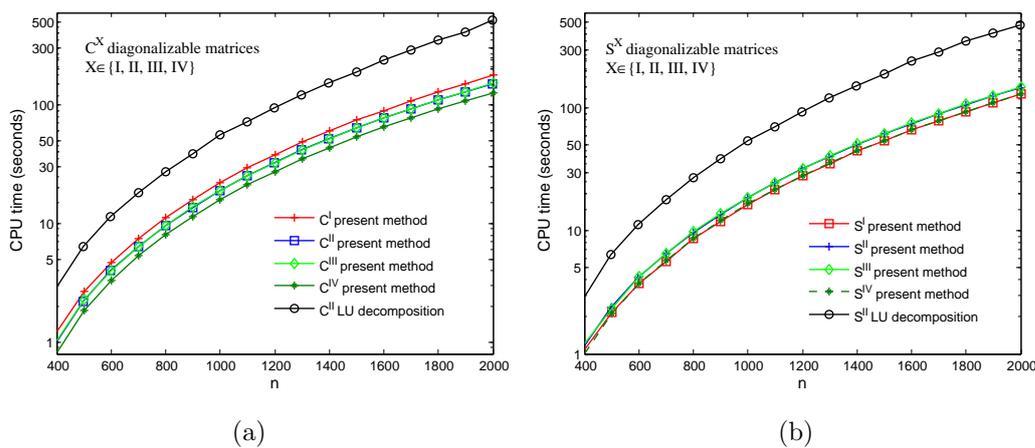


Fig. 5.3. CPU time for the inversions with respect to  $n$  of the present method and the LU decomposition for (a)  $C_n^X$ - and (b)  $S_n^X$ -diagonalizable matrices for  $X=I, II, III, IV$ .

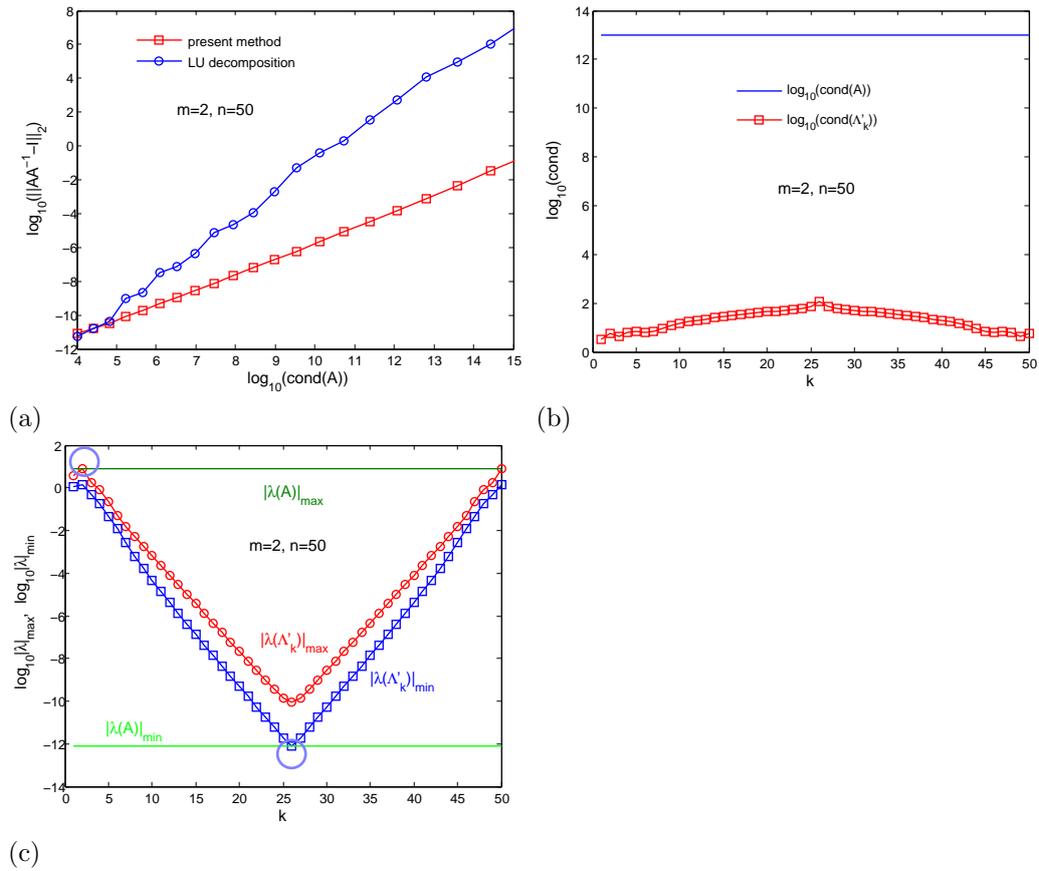


Fig. 5.4. (a) Inversion error vs condition number of a matrix  $A$  with circulant blocks for  $m=2, n=50$ , as obtained by the present method and the LU decomposition, (b) Condition number of the matrices  $\Lambda'_k$ , (c) Maximum and minimum eigenvalues of  $A$  and  $\Lambda'_k$ . The considered matrices appear in the numerical solution of wave scattering problems with the Method of Auxiliary Sources.

Table 5.1: The vectors  $\mathbf{t}$  and  $\mathbf{h}$ , which determine respectively the Toeplitz and the Hankel counterparts of the symmetric Toeplitz-plus-Hankel matrices of  $H_n^X$  matrix algebras. For  $X=I, III$  and  $X=II, IV$   $\mathbf{t}$  is the first row of the circulant and skew-circulant matrix. For  $X=I, II$  and  $X=III, IV$   $\mathbf{h}$  is the first column and last row respectively of the corresponding Hankel matrix. ( $\mathbf{x} \in \mathbb{C}^n$ )

X	$\mathbf{x}^+$	$\mathbf{x}^-$	$\mathbf{t}$	$\mathbf{h}$
I	$\mathbf{x}^+ = \frac{1}{2}(\mathbf{x} + J'_n \mathbf{x})$	$\mathbf{x}^- = \frac{1}{2}(\mathbf{x} - J'_n \mathbf{x})$	$H_n^I \mathbf{x}^+$	$H_n^I \mathbf{x}^-$
II	$\mathbf{x}^+ = \frac{1}{2}(\mathbf{x} + J_n \mathbf{x})$	$\mathbf{x}^- = \frac{1}{2}(\mathbf{x} - J_n \mathbf{x})$	$H_n^{II} \mathbf{x}^+$	$H_n^{II} \mathbf{x}^-$
III	$\mathbf{x}^+ = \frac{1}{2}(\mathbf{x} + J'_n \mathbf{x})$	$\mathbf{x}^- = \frac{1}{2}(\mathbf{x} - J'_n \mathbf{x})$	$H_n^I \mathbf{x}^+$	$H_n^I \mathbf{x}^-$
IV	$\mathbf{x}^+ = \frac{1}{2}(\mathbf{x} + J_n \mathbf{x})$	$\mathbf{x}^- = \frac{1}{2}(\mathbf{x} - J_n \mathbf{x})$	$H_n^{II} \mathbf{x}^+$	$-H_n^{II} \mathbf{x}^-$

5.2. Accuracy

We examine the accuracy of the proposed inversion method for ill-conditioned block matrices, appearing in the real-world application concerning the numerical solution of electromagnetic scattering problems by layered objects with the Method of Auxiliary Sources (MAS) [9]- [10]. Here,  $m$  is related to the number of auxiliary surfaces utilized, and  $n$  to the number of auxiliary

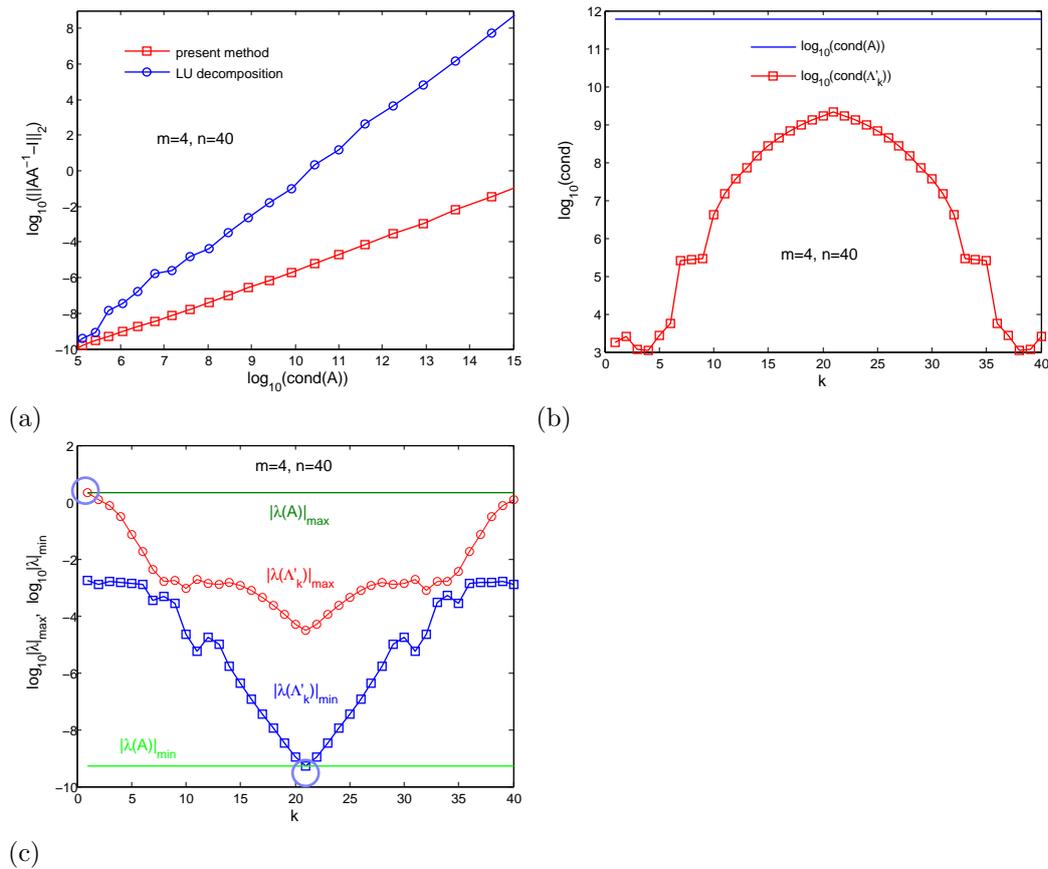


Fig. 5.5. Same as in Fig. 5.4, except for  $m=4, n=40$ .

sources considered in each auxiliary surface. Figs. 5.4a, 5.4b, and 5.4c depict respectively: (a) the inversion error  $\|AA^{-1} - I\|_2$  as a function of the condition number of a matrix  $A$  with circulant blocks with  $m=2, n=50$ , as obtained by the present method and the LU decomposition, (b) the condition number of the matrices  $\Lambda'_k$  appearing in the similarity transformation algorithm of Section 3 for a specific choice of an ill-conditioned matrix  $A$ , and (c) the maximum and minimum eigenvalues of  $A$  and  $\Lambda'_k$ . Moreover, Fig. 5.5 depicts the quantities mentioned in Fig. 5.4 but for  $m=4, n=40$ . The circles in Figs. 5.4c and 5.5c indicate the blocks  $\Lambda'_k$  containing the maximum and minimum eigenvalues of the initial matrix  $A$ .

By the statements of Figs. 5.4a and 5.5a, the achieved inversion error for ill-conditioned matrices  $A$  is significantly reduced by the application of the developed method compared with that of the LU decomposition, applied to  $A$ . This fact is justified numerically in Figs. 5.4b and 5.5b by the decrease of the conditions numbers of the blocks  $\Lambda'_k$ , which have to be inverted by the present algorithm. Hence, the problem of inverting an ill-conditioned matrix  $A$  ( $\text{cond}(A) = 10^{13}$  for  $m=2, n=50$  and  $\text{cond}(A) = 6 \cdot 10^{11}$  for  $m=4, n=40$ ) is decomposed to that of inverting the  $n$  matrices  $\Lambda'_k$ , which are no longer ill-conditioned ( $\text{cond}(\Lambda'_k) \leq 10^{2.06}$  for  $m=2, n=50$  and  $\text{cond}(\Lambda'_k) \leq 10^{9.34}$  for  $m=4, n=40$ ). This decrease of the condition numbers of matrices  $\Lambda'_k$  is further justified by the *shift of the eigenvalues* of matrix  $A$  to the eigenvalues of the blocks  $\Lambda'_k$ . More precisely, we observe from Figs. 5.4c and 5.5c that the maximum and minimum

eigenvalues of  $A$  are shifted to the blocks  $\Lambda'_2$  and  $\Lambda'_{26}$  for  $m=2$ ,  $n=50$  and to  $\Lambda'_1$  and  $\Lambda'_{21}$  for  $m=4$ ,  $n=40$  respectively.

Thus, the utilization of the developed inversion method may overcome the numerical instability problems, arising in the MAS, and hence enlarge its applicability range.

## Appendix

The Matlab function *inv\_unit\_diagon\_blocks* implements the algorithmic inversion of Section 3. This function calls the functions *dtm* and *idtm*, which implement the DTT's and the Inverse DTT's. The function *eigenmatrix* appearing in *dtm* and *idtm* generates the 16 Fourier, Hartley, Cosine, and Sine matrices, described in (2.2)-(2.5).

```
function [inv_A]=inv_unit_diagon_blocks(A,m,U,type)
% Input:  A = mxm complex block-matrix with nxn U-diagonalizable blocks
%         m = number of blocks in one block-row of A
%         U = 'F' (Fourier), 'H' (Hartley), 'C' (Cosine), 'S' (Sine)
%         type = transformation type: 1 (I), 2 (II), 3 (III), 4 (IV)
% Output: inv_A = the inverse of A
n=length(A)/m;
% Block-Diagonalization of A: creation of Lambda by using the DTT's
for i=1:m
    for j=1:m
        A_block=A((i-1)*n+1:i*n,(j-1)*n+1:j*n);
        Lambda((i-1)*n+1:i*n,(j-1)*n+1:j*n)=dtm(A_block,n,U,type);
    end
end
% Similarity Transformation: Lambda_p=P*Lambda*P';
p=1; p=horzcat(p,[1:m-1]*n+1); p_new=p; for j=2:n, p_new=p_new+1;
p=horzcat(p,p_new); end Lambda_p=Lambda(p,:);
Lambda_p=Lambda_p(:,p);
% Computation of the inverse inv_Lambda_p of Lambda_p;
for i=1:n
    Lambda_p_block=Lambda_p((i-1)*m+1:i*m,(i-1)*m+1:i*m);
    inv_Lambda_p((i-1)*m+1:i*m,(i-1)*m+1:i*m)=inv(Lambda_p_block);
end
% Inverse Similarity Transformation: inv_Lambda=P'*inv_Lambda_p*P;
p=1; p=horzcat(p,[1:n-1]*m+1); p_new=p; for j=2:m, p_new=p_new+1;
p=horzcat(p,p_new); end inv_Lambda=inv_Lambda_p(p,:);
inv_Lambda=inv_Lambda(:,p);
% Block-Diagonalization of inv_A by imposing the IDTT's in inv_Lambda
for i=1:m
    for j=1:m
        inv_Lambda_block=inv_Lambda((i-1)*n+1:i*n,(j-1)*n+1:j*n);
        if U=='F'
            inv_A_vector=idtm(inv_Lambda_block,n,U,type);
            inv_A_block=gallery('circul',inv_A_vector);
            if (type==2) | (type==4)
```

```

        inv_A_block=triu(inv_A_block)-tril(inv_A_block,-1);
    end
    else if (U=='H') | (U=='C') | (U=='S')
        inv_A_block=idtt(inv_Lambda_block,n,U,type);
    end
    end
    inv_A((i-1)*n+1:i*n,(j-1)*n+1:j*n)=inv_A_block;
end
end
end
function out=dtm(X,n,U,type)
if U=='F'
    if (type==1) | (type==3)
        T=eigenmatrix(n,U,1); out=sqrt(n)*diag(T*X(1,:).');
    end
    if (type==2) | (type==4)
        T=eigenmatrix(n,U,1);
        W=diag(exp(-i*pi*[0:n-1]/n)); Wvec=conj(W(1,:)); WXvec=Wvec*X;
        help=WXvec*diag(exp(-i*pi*[0:n-1]/n)); out=sqrt(n)*diag(T*help. ');
    end
end
end
if (U=='H') | (U=='C') | (U=='S'), T=eigenmatrix(n,U,type); out=T'*X*T; end
end
function out=idtt(X,n,U,type)
T=eigenmatrix(n,U,type);
if U=='F', Tvec=T(1,:); TXvec=Tvec*X; out=TXvec*T';
else if (U=='H') | (U=='C') | (U=='S'), out=T*X*T'; end
end
end

```

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